# National Tsing Hua University Department of Computer Science CS1356 -- Introduction to Information Engineering Mid-term Examination, Fall 2009 

1. (20 pt) Consider an 8-bit binary number system with 2 's complement for negative numbers.
(a) What is 37 (base 10) represented in this system?
(b) What is -48 (base 10) represented in this system?
(c) What is 5A (base 16) represented in this system?
(d) What is the decimal number of 00101101 (base 2)?
(e) Give an example that can cause overflow by using this number representation.

Ans:
(a) 00100101
(b) 11010000
(c) 01011010
(d) 45
(e) There are two cases that may incur overflow:

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Positive + Positive (e.g. 127 + 3) or Negative + Negative ( e.g. -128 + (-1))
```

2. (12 pt) Consider an 8-bit floating point number whose bit pattern is 1 bit for sign, 3 bits for exponent, and 4 bits for mantissa in order. The exponent uses excess notation ranged from -4 to 3 , and the mantissa is normalized so that it is larger than or equal to $1 / 2$ and smaller than 1 .
(a) What is the binary representation of $21 / 16$ ?
(b) What is the floating point representation mentioned above for $21 / 16$ ?
(c) What is the truncation error of this floating point representation?

Ans:
(a) 1.0101
(b) $1.0101=0.10101 \times 2^{1}$ excess $=1+4=5=101$ => 01011010 (the least significant bit is truncated)
(c) The truncation error is $0.0625_{d}\left(=0.00001_{b}\right)$
3. (6pt) A Boolean function $F(x, y)$ is defined as in the truth table. Design a circuit that uses AND, OR, and NOT gates to implement the function $F(x, y)$ and justify your answer.

Ans:


| $x$ | $y$ | $F(x, y)$ |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 0 | 1 | 1 |
| 1 | 0 | 0 |
| 1 | 1 | 0 |

4. (12 pt) Suppose Alice wants to send an image file of size $\mathbf{4 M}$ bytes to Bob. The network has a transfer rate of $\mathbf{2}^{\mathbf{2 0}} \mathbf{~ b p s}$. However, for every $\mathbf{1 0 2 4}$ bits, there will be one bit error.
(a) Suppose an odd parity bit is added to every byte to detect the error. How large the file will become after the error detection bits are added?
(b) How long does it take to transfer this encoded image? (You can assume the encoded image is of size $X$ bytes and express your answer using $X$ if you cannot make sure your answer in 4(a) is correct.)
(c) After getting the entire encoded image, Bob knows which parts of file have errors. He sends a list showing errors to Alice, and asks Alice to resend those erroneous parts. How many bytes does Alice need to resend? Suppose this time, no error detection code is added and no error occurs.
Ans:
(a) $4 \mathrm{Mbytes}+4 \mathrm{Mbits}=4 \mathrm{Mbytes}+0.5 \mathrm{Mbytes}=4.5 \mathrm{Mbytes}($ or 4.5 MB$)$
(b) 4.5 M byte $/ 1 \mathrm{Mbps}=36 \mathrm{M}$ bits/ $1 \mathrm{Mbps}=36$ seconds (or $8 \mathrm{X} / \mathrm{M}$ seconds)
(c) For every 1024 bits, there will be one bit error. However, when there is one bit error, the entire byte need be resent, because Bob cannot know which bit is erroneous. (He can only know which "byte" has error.) Therefore, for every 1024 bits, there are 8 bits (1 byte) need be resent. Since the entire encoded image is 4.5 MB , the number of data to be resent is $4.5 \mathrm{M} /(1024 / 8)=36 \mathrm{~KB} \quad($ or $\mathrm{X} / 128)$
5. ( 5 pt ) Give one advantage for using the stored-program concept in designing computers and explain your reason.
6. ( 10 pt ) Using the machine language described in the language description table given at the end, write a sequence of instructions that will find the two's complement of the value stored at address AO. The result is stored in register 4. (Hint: You could use XOR to find the complements.)
Ans: 10A0
21FF
2201
9301
5423
7. (15 pt) Using the machine language described in the language description table given at the end, write a sequence of instructions to implement the following C statement. Assume that the variable $A$ is stored at address D0 with an initial value of F9. (Hint: Use shift and addition instructions to implement * and /. Note that the numbers have signs.)

$$
A=(A * 5) / 2 ;
$$

What is the final value in address DO assuming a two's complement representation?

Ans:
Idea: 1, Calculate A*5 by Branch operation Address AO ~ BO
2. Calculate $A^{*} 5 / 2$ by Rotate Address BA
3. Sign extension Address B2~BC

Address
AO 2005
A2 2101
A4 2200
A6 13D0
A8 2400
AA 5434
AC 5221
AE B2B2
BO BOAA
B2 21FE
B4 8414
B6 2280
B8 8224
BA A401
BC 7424
BE 34DO
8. (10 pt) The following C program adds 3 numbers in an array A.

```
int main(){
    int A[3]={3,4,2}, B=0, i;
    for(i=0; i<3; i++) B=B+A[i];
}
```

Suppose the program is translated to the machine program, written in the machine language described in the language description table given at the end, and is stored in main memory beginning at address $\mathbf{3 0}$ (hexadecimal). The data type "int" is 8 bits in the machine. Array A is stored in main memory 00
to $\mathbf{0 2}$. Variable $\mathbf{B}$ is stored at $\mathbf{0 3}$. Their initial values are set before the program is executed.
(a) The translated program is partially given in the left. To make the program

B2 $\qquad$
$\qquad$ execute correctly, what should be in the blank places?
(b) What happen to the program if the statement " 3239 " is erroneously entered as " 3238 "?

Ans:
(a)

(b)

If 3239 is replaced with 3238 , it will change the OP code which is in the " 38 " address of memory. Then the program logic will be wrong.
9. (10 pt) Suppose in a fast food restaurant, a hamburger menu consists of a hamburger, a drink, and French fries. Their regular prices are $\$ 60, \$ 20$, and $\$ 20$, respectively. Since this is your birthday, you are able to negotiate a special deal for the hamburger at $\$ 15$. Therefore, you get a saving of

$$
\text { saving }=1-\frac{\text { price of a discount hamburger menu }}{\text { price of a regular hamburger menu }}=1-\frac{15+20+20}{60+20+20}=45 \%
$$

This saving can be explained using Amdahl's law, which states that the speedup using $n$ processors to execute a program with an $f$ portion $(0<f \leq 1)$ that cannot be parallelized is

$$
\text { speedup }=\frac{\text { execution time uses one processor }}{\text { execution time uses } n \text { processors }}=\frac{t_{s}}{t_{p}}=\frac{t_{s}}{f \cdot t_{s}+(1-f) \cdot t_{s} / n}=\frac{n}{1+(n-1) f}
$$

(a) What are $f$ and $n$ in the hamburger case?
$\mathrm{N}=60 / 15=4$ (In Amdahl's law, the n represents the number of processors. In this case, the cost of Hamburger is reduced with 15.)
$F=20+20 / 60+20+20=0.4=40 \%$
(b) If you are able to use your charm to negotiate even lower prices for the hamburger, what is the maximum saving that you can expect?
$1-(20+20) /(60+20+20)=60 \%$

Language Description Table

| Opcode | Operand | Description |
| :---: | :---: | :---: |
| 1 | RXY | LOAD the register R with the bit pattern found in the memory cell whose address is XY . |
| 2 | RXY | LOAD the register R with the bit pattern XY . |
| 3 | RST | STORE the bit pattern found in register $R$ in the memory cell whose address is XY . |
| 4 | ORS | MOVE the bit pattern found in register R to register S . |
| 5 | RST | ADD the bit patterns in registers $S$ and $T$ as though they were two's complement representations and leave the result in register $R$. |
| 6 | RST | ADD the bit patterns in registers $S$ and $T$ as though they represented values in floating point notation and leave the floating-point result in register R. |
| 7 | RST | OR the bit patterns in registers $S$ and $T$ and place the result in register R. |
| 8 | RST | AND the bit patterns in registers $S$ and $T$ and place the result in register $R$. |
| 9 | RST | EXCLUSIVE OR the bit patterns in registers $S$ and $T$ and place the result in register R. |
| A | ROX | ROTATE the bit pattern in register R one bit to the right X times. Each time place the bit that started at the low-order end at the high-order end. |
| B | RXY | JUMP to the instruction located in the memory cell at address XY if the bit pattern in register $R$ is equal to the bit pattern in register number 0 . Otherwise, continue with the normal sequence of execution. |
| C | 000 | HALT execution. |

Statistics of midterm: Average: 61.93
Histogram


